

# Combinatorial Games Crash Course

Kyle Burke

Plymouth State

April 5, 2019

# Combinatorial Games

Combinatorial Games:

- ▶ Two competing players ("Left", "Right") alternate turns
- ▶ No randomness
- ▶ No hidden information
- ▶ If you can't move, you lose. ("Normal" play)

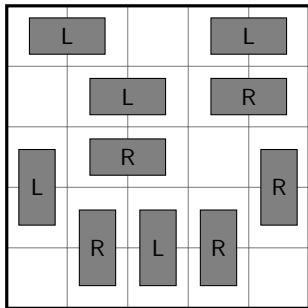
We can describe games based on the moves available to each player.

# Cram

Cram:

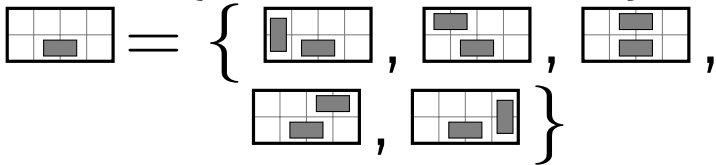
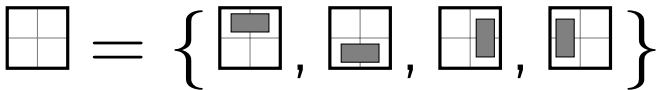
- ▶ Start with an empty checkerboard.
- ▶ Turn: play a domino on any two empty adjacent squares.

Quick playthrough:

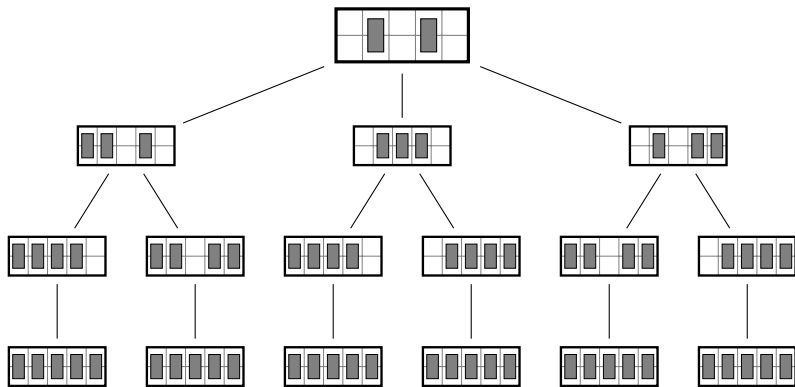


# Options

We can abstractly describe each position by the *options*, positions that can be moved to.

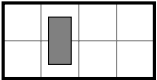




# Game Trees



# Adding Games

Games break into pieces, which we want to treat independently.

We want:  to act the same as  + 

# Adding Games

So,  $G + H$  means on their turn, the player chooses an option from either  $G$  or  $H$  (not both).

Rigor:  $G = \{G_1, \dots, G_n\}$  and  $H = \{H_1, \dots, H_m\}$ , then

$$G + H = \{G_1 + H, \dots, G_n + H, H_1 + G, \dots, H_m + G\}$$

Most of our definitions are going to revolve around how addition works.

## Another game: Nim

Nim:

- ▶ Heap(s) of tokens/sticks
- ▶ Turn: remove some sticks from any one heap.

Position (as list): [1, 2, 3]

Options of [1, 2, 3]:

- ▶ (Take 1 from first) [2, 3],
- ▶ (Take 1 from second) [1, 1, 3],
- ▶ [1, 3],
- ▶ [1, 2, 2],
- ▶ [1, 2, 1],
- ▶ [1, 2]

$[1, 2, 3] = \{ [2, 3], [1, 1, 3], [1, 3], [1, 2, 2], [1, 2, 1], [1, 2] \}$



## Zero and \*

{ } seems important. Let's give it a name.

$$\{ \} := 0$$

Let's do {0} too:

$$\{0\} := *$$

# Equivalence

What happens when we add 0 to another game,  $G$ ?

Same as  $G$

Equivalence in general?  $G = H$ ?

$G = H \Leftrightarrow \forall X : G + X$  has the same outcome as  $H + X$

What do we mean by *outcome*?

# Outcome Classes

Outcome class: set of games where the same player always has a winning strategy.

Who wins on 0? ("Next" or "Previous" player?) A: Previous

Let's call the set of these positions  $\mathcal{P}$ , "Zero".

Bonus:  $\forall X \in \mathcal{P} : X = 0$

# Outcome Classes

Who wins on  $*$ ? A: Next player!

Let's call these positions  $\mathcal{N}$ , "Fuzzy".

What's true of all single Nim heaps with at least one stick?

They're all in fuzzy!

Not all fuzzy games are equal to some single value, so they're...  
"fuzzy".

## Fuzzy + Fuzzy

$$\square \in \mathcal{N}$$

What about  $\square + \square$ ?

Sum  $\in \mathcal{P}$ , so it equals 0.

## Fuzzy + Fuzzy

$$\begin{array}{|c|} \hline \\ \hline \\ \hline \end{array} \in \mathcal{N}, \quad \begin{array}{|c|c|c|c|} \hline \\ \hline \\ \hline \end{array} \in \mathcal{N}$$

What about  $\begin{array}{|c|} \hline \\ \hline \\ \hline \end{array} + \begin{array}{|c|c|c|c|} \hline \\ \hline \\ \hline \end{array}$ ?

Still  $\mathcal{N}$

Winning Move:  $\begin{array}{|c|} \hline \\ \hline \\ \hline \end{array} + \begin{array}{|c|c|c|c|} \hline \\ \hline \\ \hline \end{array}$

Outcome class defined by *existence* of a strategy. Not every move will be winning.

## Different Fuzzy Values: Cram

$$\begin{aligned} \boxed{\begin{array}{|c|c|c|c|} \hline & & & \\ \hline \end{array}} &= \left\{ \boxed{\begin{array}{|c|c|c|c|} \hline \blacksquare & & & \\ \hline \end{array}}, \boxed{\begin{array}{|c|c|c|c|} \hline & \blacksquare & & \\ \hline \end{array}}, \boxed{\begin{array}{|c|c|c|c|} \hline & & & \blacksquare \\ \hline \end{array}} \right\} \\ &= \{*, 0, *\} = \{0, *\} \end{aligned}$$

Can win, but also has more power...

Let's call this \*2.

## Different Fuzzy Values: Nim

What's the value of the single Nim heap: [2]?

$$\begin{aligned}[2] &= \{ [], [1] \} \\ &= \{ 0, * \} \\ &= *2\end{aligned}$$

What's the value of [3]?

$$\begin{aligned}[3] &= \{ [], [1], [2] \} \\ &= \{ 0, *, *2 \} \\ &= *3\end{aligned}$$

True for all single Nim heaps!  $[k] = *k$



# Numbers!

More generally,  $*k = \{ *0, *1, *2, *3, \dots, *(k-1) \}$

" $*k$  has all numbers 0 through  $k-1$ , but not  $k$ "

Note: it doesn't matter whether there are numbers above  $k$ .

"Minimal Excluded Value" (mex)

$$\text{mex}(S \subset (\mathbb{N} \cup \{0\})) = \min_{x \notin S} (\mathbb{N})$$

# Cram Nimbers

$$\square = \{ \} = 0$$

$$\square \square = \{0\} = *$$

$$\square \square \square = \{0,0\} = \{0\} = *$$

$$\square \square \square \square = \{*,0,*\} = *2$$

$$\square \square \square \square \square = \{ \square \square \square \square \square \square, \square \square \square \square \square \square, \square \square \square \square \square \square, \square \square \square \square \square \square \} = \{ *, *, *, * \} = 0$$

$$\square \square \square \square \square \square = \{ *2, *, 0 \} = *3$$

$$\square \square \square \square \square \square \square = \{ 0, *2, 0 \} = *$$

$$\square \square \square \square \square \square \square \square = \{ *3, 0, *3, 0 \} = *$$

# Adding Numbers

Bonus: summing numbers is easy!

$$*k + *m = *(k \oplus m)$$

$$\begin{aligned} *6 + *5 &= *(110 \oplus 101) \\ &\quad 110 \\ &\quad \oplus 101 \\ &= \quad 011 \\ &= \quad *3 \end{aligned}$$

- ▶  $*8 + *7 = *15$
- ▶  $* + * = 0$ .  $*k + *k = 0$
- ▶  $*10 + *13 = *7$

# Cram Nimbers

$$\boxed{\phantom{00}} = *$$

$$\boxed{\phantom{000}} = *$$

$$\boxed{\phantom{0000}} = *2$$

$$\boxed{\phantom{00000}} = \{*, *, *, *\} = 0$$

$$\boxed{\phantom{000000}} = \{*2, *, 0\} = *3$$

$$\boxed{\phantom{0000000}} = \{0, *2, 0\} = *$$

$$\boxed{\phantom{00000000}} = \{*3, 0, *3, 0\} = *$$

# Impartial vs Partisan

So far both players have always had the same moves. ("Impartial games")

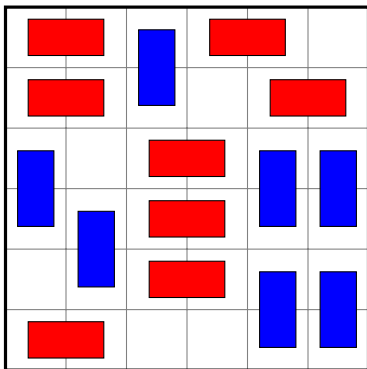
Nimbers cover all values for impartial games.

Different: "Partisan games". Player identity matters!

Left vs Right: Blue vs Red.

- ▶ New outcome classes!
- ▶ New representation!
- ▶ New values!

# Domineering



# Partisan Outcome Classes

Previously...

- ▶  $\mathcal{P}$ : Previous player wins, "Zero"
- ▶  $\mathcal{N}$ : Next player wins, "Fuzzy"

New possibilities:

- ▶  $\mathcal{L}$ : Left wins, "Positive"
- ▶  $\mathcal{R}$ : Right wins, "Negative"

Notation:  $o(G) =$  outcome class of  $G$

# Domineering Outcomes

$$\square \in \mathcal{P}$$

$$\begin{array}{|c|c|} \hline & \\ \hline & \\ \hline \end{array} \in \mathcal{N}$$

$$\begin{array}{|c|} \hline \\ \hline \\ \hline \end{array} \in \mathcal{L}$$

$$\begin{array}{|c|c|} \hline & \\ \hline & \\ \hline \end{array} \in \mathcal{R}$$

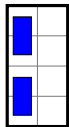


# 1 and 2



= 1

"One move for Left"



= 2

"Two moves for Left"

# Partisan Representations

{ Left's options | Right's options }

- ▶  $0 = \{ \mid \}$
- ▶  $* = \{0 \mid 0\}$
- ▶  $*2 = \{0, * \mid 0, *\}$
- ▶ etc...

What happens if the two sides aren't the same?

## Partisan Representations

{ Left's options | Right's options }

$$\begin{array}{|c|c|} \hline \color{blue}{\blacksquare} & \\ \hline \color{blue}{\blacksquare} & \\ \hline \end{array} = 1 = \{0 \mid \}$$

$$\begin{array}{|c|c|} \hline \color{blue}{\blacksquare} & \\ \hline \color{blue}{\blacksquare} & \\ \hline \end{array} = 2 = \{1, 0 \mid \}$$

We'd really like to drop that 0 inside of 2...

Intuition:  $0 < 1$ , so we can remove the 0 option, because Left will always choose 1 instead. How can we define  $<$ ?

# Comparing Games

$G < H$  if Left never prefers  $G$  to  $H$

Never? In the context of other games,  $X$

- ▶  $G \leq H$  if Left wins on  $G + X \Rightarrow$  Left wins on  $H + X$
- ▶ Recall:  $G = H$  if  $\forall X : o(G + X) = o(H + X)$
- ▶ So:  $G < H$  if  $G \leq H$  and  $\exists X : o(G + X) \neq \mathcal{L} = o(H + X)$

$0 < 1$  and  $* < 1$

- ▶  $G \parallel H$  if  $G \not\leq H$  and  $G \not\leq H$

$0 \parallel *$  and  $0 \parallel * 2$

## Dominated Moves & Integers

$$\{0, 1 \mid \dots\} = \{1 \mid \dots\}$$

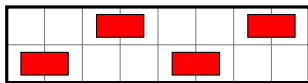
$$\{\dots \mid *, 1\} = \{\dots \mid *\}$$

$$\{1 \mid \} = 2$$

$\forall k \in \mathbb{N}$ :

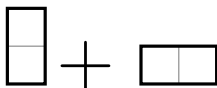
▶  $\{k \mid \} = k + 1$

▶  $\{ \mid -k\} = -k - 1$



## Negatives

What's the outcome class of  $1 + -1$ ?



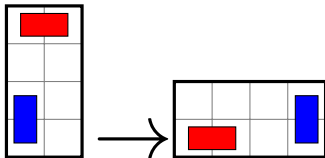
$$1 + -1 \in \mathcal{P}$$

In general,

$$\begin{aligned} -G &= -\{L_1, \dots, L_n \mid R_1, \dots, R_m\} \\ &= \{-R_1, \dots, -R_m \mid -L_1, \dots, -L_n\} \\ -(k+1) &= -\{k \mid \} \\ &= \{ \mid -k \} \\ &= -k - 1 \end{aligned}$$

# Domineering Negatives

In Domineering, just reflect and switch colors!



# Domineering Halves

What's this sum?

$$\begin{array}{|c|} \hline \square \\ \hline \square \\ \hline \end{array} + \begin{array}{|c|c|c|} \hline \square & \square & \blacksquare \\ \hline \blacksquare & \square & \square \\ \hline \end{array} + \begin{array}{|c|c|c|} \hline \square & \square & \blacksquare \\ \hline \blacksquare & \square & \square \\ \hline \end{array} = 0$$

So...  $\begin{array}{|c|c|c|} \hline \square & \square & \blacksquare \\ \hline \blacksquare & \square & \square \\ \hline \end{array} + \begin{array}{|c|c|c|} \hline \square & \square & \blacksquare \\ \hline \blacksquare & \square & \square \\ \hline \end{array} = -1$

What about this?

$$\begin{array}{|c|} \hline \square \\ \hline \square \\ \hline \end{array} + \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \blacksquare & \blacksquare & \square \\ \hline \end{array} + \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \blacksquare & \blacksquare & \square \\ \hline \end{array} = 0$$

Same thing!



## Domineering Halves

What about this?

$$\begin{array}{|c|} \hline \square \\ \hline \square \\ \hline \end{array} + \begin{array}{|c|c|c|} \hline \square & \square & \blacksquare \\ \hline \blacksquare & \square & \square \\ \hline \end{array} + \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \blacksquare & \blacksquare & \square \\ \hline \end{array} \neq 0$$

$$\text{So, } \begin{array}{|c|c|c|} \hline \square & \square & \blacksquare \\ \hline \blacksquare & \square & \square \\ \hline \end{array} + \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \blacksquare & \blacksquare & \square \\ \hline \end{array} \neq -1$$

$$\begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \blacksquare & \blacksquare & \square \\ \hline \end{array} = \begin{array}{|c|c|c|} \hline \square & \square & \blacksquare \\ \hline \blacksquare & \square & \square \\ \hline \end{array} ?$$

$$\begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \blacksquare & \blacksquare & \square \\ \hline \end{array} + \begin{array}{|c|} \hline \square & \blacksquare \\ \hline \square & \square \\ \hline \blacksquare & \square \\ \hline \end{array} = 0?$$

Only one of these is  $-1/2$ ...

# Domineering Halves

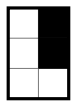
Look at the outcome classes!

$$\begin{array}{|c|c|} \hline \blacksquare & \blacksquare \\ \hline \square & \blacksquare \\ \hline \square & \blacksquare \\ \hline \square & \square \\ \hline \end{array} \in \mathcal{L} \quad (= 1/2)$$

$$\begin{array}{|c|c|} \hline \blacksquare & \blacksquare \\ \hline \square & \square \\ \hline \square & \square \\ \hline \blacksquare & \square \\ \hline \end{array} \in \mathcal{N} \quad (= 1/2 + *)$$

Two times either one is 1

# Simplest Numbers


$$= \{0 \mid 1\} = 1/2$$

Simplest numbers!

If  $m < n$  (both numbers), then  $\{m \mid n\} = k$ , where:

- ▶ If there is any integer strictly between  $m$  and  $n$ , (i.e.,  $k \in (m, n)$ ) then  $k$  is the one with smallest absolute value.
- ▶ Otherwise,

$$k = \frac{a}{2^b}$$

where  $m < k < n$  and  $b > 0$  is minimal.

# Simplest Numbers

Examples!

- ▶  $\{0 \mid 5\} = 1$
- ▶  $\{-99 \mid 5\} = 0$
- ▶  $\{99 \mid 103\} = 100$
- ▶  $\{0 \mid 1\} = 1/2$
- ▶  $\{3/4 \mid 1\} = 7/8$

What about this?

$$\begin{array}{|c|c|} \hline & \\ \hline & \\ \hline \end{array} = \{ \mathbf{1} \mid - \mathbf{1} \}$$

No simplest number here!

# Switches

$$\boxed{\begin{array}{|c|c|} \hline & \\ \hline & \\ \hline \end{array}} = \{1 \mid -1\} \neq 0$$

In general,  $\pm k = \{k \mid -k\}$

- ▶ Not surreal
- ▶ "Hot" games
- ▶  $* = \pm 0$
- ▶  $\{6 \mid 0\} = 3 + \{3 \mid -3\} = 3 \pm 3$
- ▶ Sums possible! E.g.  $2 \pm 6 \pm 3 \pm 2$

# Infinitesimals

$\uparrow = \{0 \mid *\}$

- ▶  $\in \mathcal{L}$
- ▶ Smaller than any positive number!
- ▶ "Dicot" games
- ▶ Even smaller: Tinys and Minys!

# So much more!

"Math":

- ▶ Which values from which rulesets?
- ▶ Misere Play, Scoring Games
- ▶ Other Sums

"Computer Science":

- ▶ Computational Complexity
- ▶ AI

Thank you!